

Divide and Conquer Matrix mulitplication **Design and Analysis**

of Algorithms I





EXAMPLE (N=2) $\binom{a \ b}{(e \ f)} = \binom{ae+bg}{(e+bg)} = \binom{af+bh}{(e+bg)}$

Nextcore AI -Gopal Shangari



What is the asymptotic running time of the straightforward iterative algorithm for matrix multiplication?

 $\bigcirc \theta(n \log n) \\ \bigcirc \theta(n^2)$

 $> \bigcirc \theta(n^3)$

NEXTCORE AI

 $\bigcirc \theta(n^4)$



THE DIVIDE AND CONQUER PARADIGM

- 1. DIVIDE into smaller subproblems
- 2. CONQUER subproblems recursively.
- 3. COMBINE solutions of subproblems into one for the original problem.



$\frac{\text{Idea}}{\text{Write } \forall = \begin{pmatrix} e & e \\ c & e \end{pmatrix}_{\text{and}} \forall = \begin{pmatrix} e & e \\ c & e \end{pmatrix}$

[where A through H are all n/2 by n/2 matrices]

Then : (you check)

$$Y = \begin{pmatrix} AE + TSG & AF + BH \\ CE + DG & CF + DH \end{pmatrix}$$



RECURSIVE ALGORITHM #1

<u>Step 1</u> : recursively compute the 8 necessary products.

<u>Step 2</u> : do the necessary additions $(\theta(n^2) \ time)$

<u>Fact</u> : runtime is $\theta(n^3)$ [follows from the master method]



Strassen's Algorithm (1969)

<u>Step 1</u> : recursively compute only 7 (cleverly chosen) products

<u>Step 2</u> : do the necessary (clever) additions + subtractions (still $\theta(n^2)$ time)

<u>Fact</u> : better than cubic time!

[see Master Method lecture]



Y= (EF) $\chi = (c)$ The Details <u>The Seven Products</u>: $P_1 = A(F-H)$, $P_2 = (A+B)H$, $P_{7} = ((+D)E, P_{4} = D(G-E), P_{5} = (A+D)(E+H).$ $P_{1} = (9 - D)(6 + H), P_{7} = (A - C)(E + F)$ Claim: X.Y = (ce + DG CF + DH) = (PS + P4 - P2 + PG P1 + P2) (P3 + P4 (P1 + P5 - P3 - P3) Proof: AE+AH+DE+DH+DG-DE-AH-BH Q.E.D (remains + SG + BH - DG - DH = AE + BG <u>Question</u> : where did this come from ? open!)