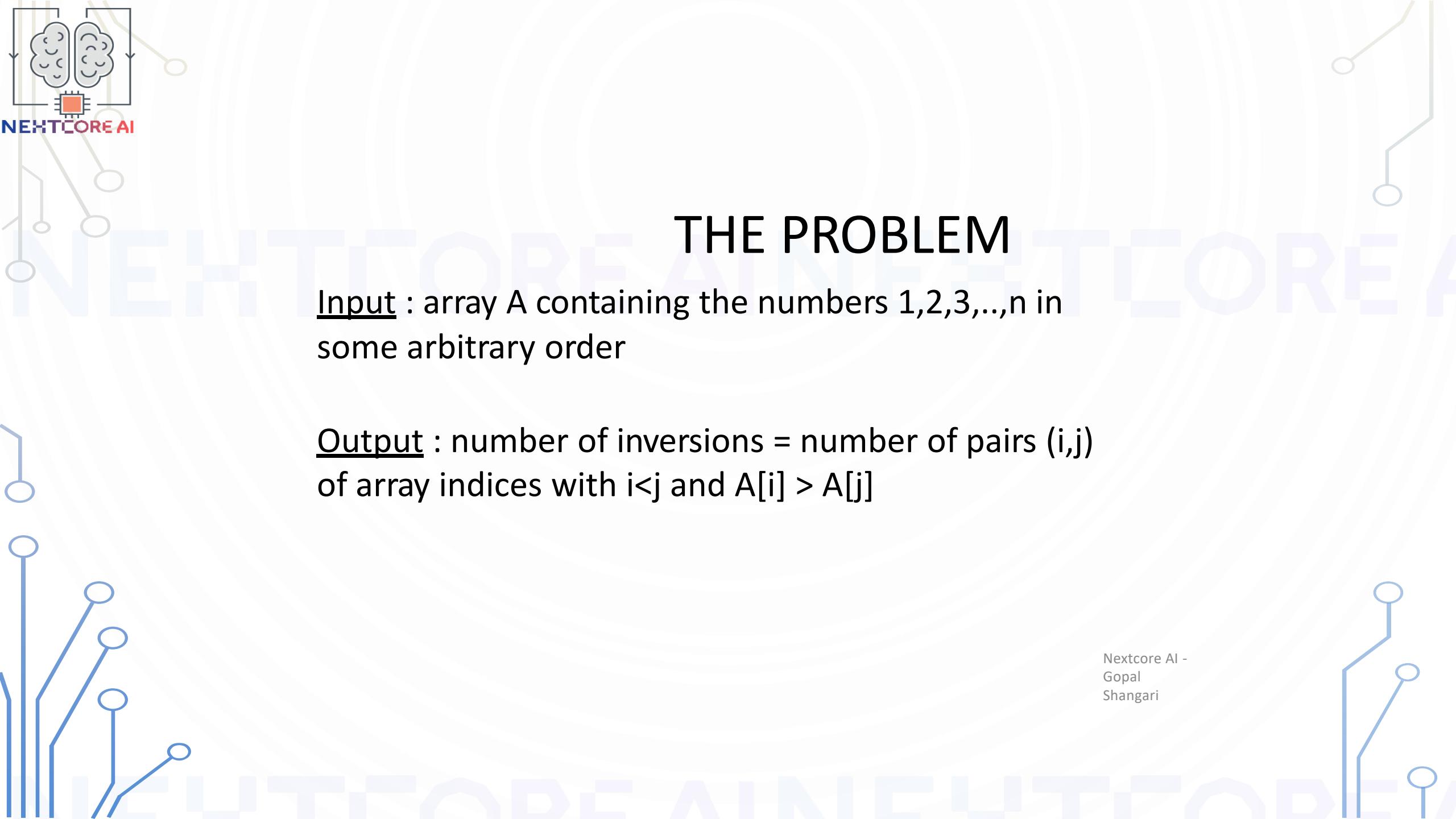


Design and Analysis
of Algorithms I

Divide and Conquer

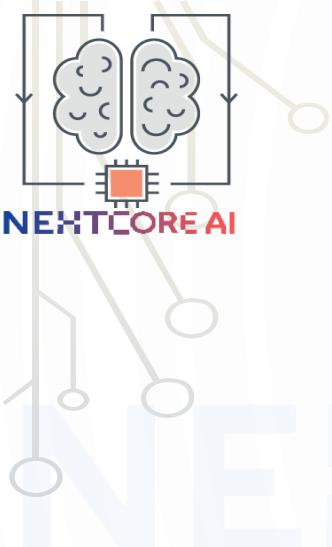
Counting Inversions I



THE PROBLEM

Input : array A containing the numbers 1,2,3,..,n in some arbitrary order

Output : number of inversions = number of pairs (i,j) of array indices with $i < j$ and $A[i] > A[j]$



Examples and Motivation

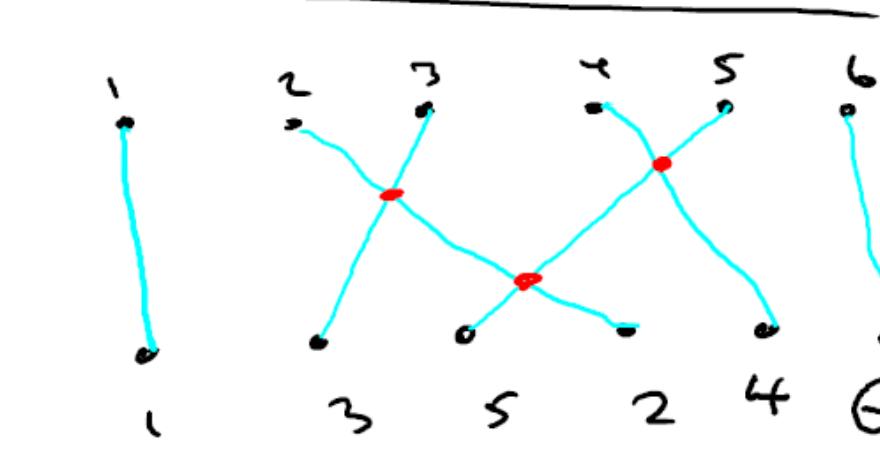
Example

(1, 3, 5, 2, 4, 6)

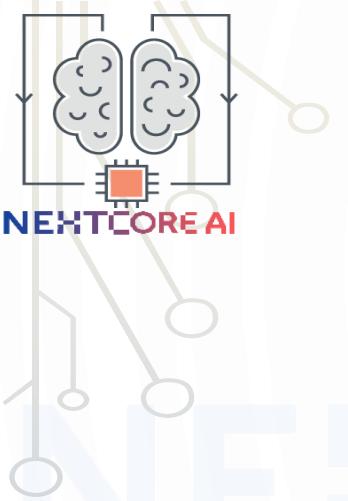
Inversions :

(3,2), (5,2), (5,4)

Motivation : numerical
similarity measure
between two ranked lists eg: for collaborative filtering



Gopal
Shangari



NEXTCORE AI

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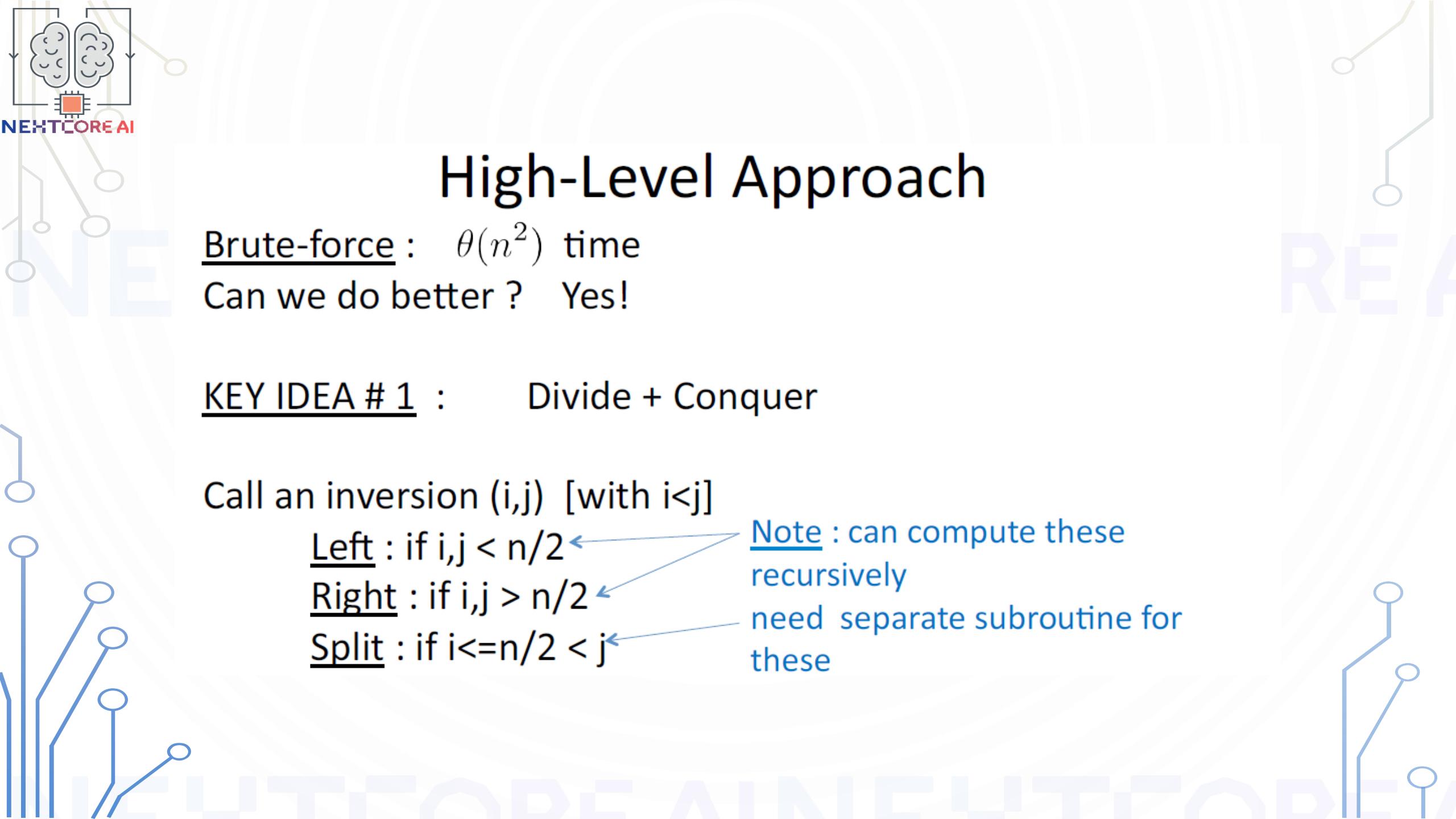
What is the largest possible number of inversions that a 6-element array can have?

15 In general, $\binom{n}{2} = n(n - 1)/2$

21

36

64



High-Level Approach

Brute-force : $\theta(n^2)$ time

Can we do better ? Yes!

KEY IDEA # 1 : Divide + Conquer

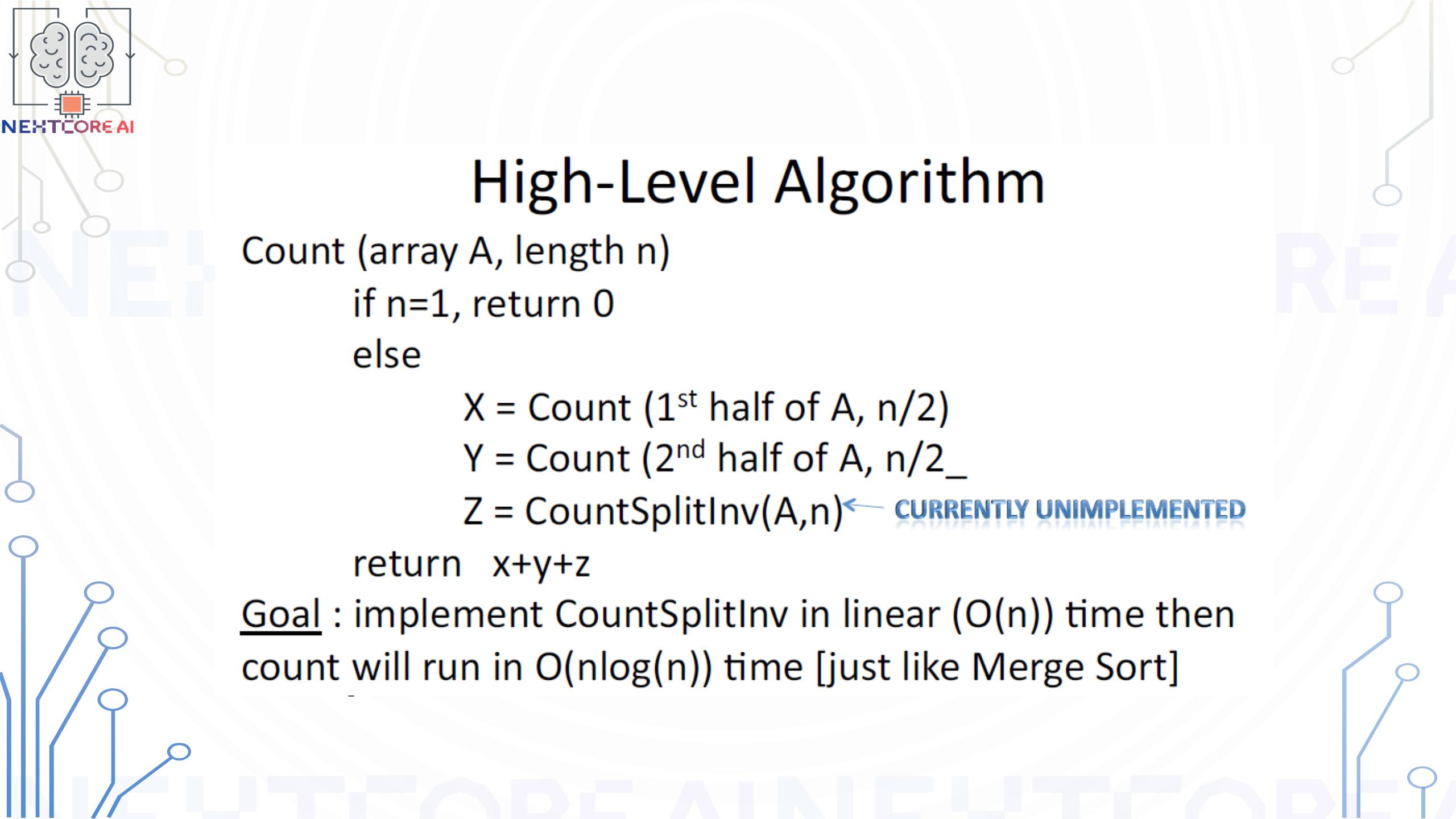
Call an inversion (i,j) [with $i < j$]

Left : if $i, j < n/2$

Right : if $i, j > n/2$

Split : if $i \leq n/2 < j$

Note : can compute these
recursively
need separate subroutine for
these



High-Level Algorithm

Count (array A, length n)

 if n=1, return 0

 else

 X = Count (1st half of A, n/2)

 Y = Count (2nd half of A, n/2)

 Z = CountSplitInv(A,n) ← CURRENTLY UNIMPLEMENTED

 return x+y+z

Goal : implement CountSplitInv in linear ($O(n)$) time then
count will run in $O(n \log(n))$ time [just like Merge Sort]